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Discrete Applied Mathematics 113 (2001) 143–166

DISCRETE
APPLIED
MATHEMATICS

Diameter determination on restricted graph families [☆]

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Received 25 August 1998; revised 1 June 2000; accepted 12 June 2000

Abstract

Determining the diameter of a graph is a fundamental graph operation, yet no efficient (i.e. linear or quadratic time) algorithm is known. In this paper, we examine the diameter problem on chordal graphs and AT-free graphs and show that a very simple (linear time) 2-sweep LexBFS algorithm identifies a vertex of maximum eccentricity unless the given graph has a specified induced subgraph (it was previously known that a single LexBFS algorithm is guaranteed to end at a vertex that is within 1 of the diameter for chordal graphs and AT-free graphs). As a consequence of the forbidden induced subgraph result on chordal graphs, our algorithm is guaranteed to work optimally for directed path graphs (it was previously known that a single LexBFS algorithm is guaranteed to work optimally for interval graphs). © 2001 Elsevier Science B.V. All rights reserved.

1. Introduction

Recently considerable attention has been given to the problem of developing fast and simple algorithms for various classical graph problems. The motivation for such algorithms stems from our need to solve these problems on very large input graphs, thus the algorithms must be not only fast, but also easily implementable. Determining a graph diameter is a classical and well-known problem.

For arbitrary graphs (with n vertices and m edges), as well as for various restricted graph families, the current fastest algorithm for this problem achieves the time bound of $O(nm)$ (see for example [24]) which is too slow to be practical for very large graphs.

[☆] An extended abstract of these results has been presented at WG'98 [7].

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This naive algorithm performs breadth first searches (BFSs) from each vertex of the graph and actually finds the whole distance matrix of the graph. For dense graphs, the best result known is by Seidel [28], who showed that the distance matrix (and hence the diameter) of a graph can be computed in $O(M(n)\log n)$ time where $M(n)$ denotes the time complexity for matrix multiplication involving small integers only. (Currently, $M(n)$ is known to be $O(n^{2.376})$ [6].) Note also that in a recent paper [1], the authors solve the all pairs shortest-path problem with an additive error at most 2 without matrix multiplication in $O(n^{2.5}\sqrt{\log n})$ time. They obtain also a ratio $\frac{2}{3}$ approximation to the diameter in time $O(m\sqrt{n\log n} + n^2\log n)$.

Algorithms of complexity $O(n^2)$ for computing the distance matrix (and hence the diameter) in different particular graph classes were presented in [10,2,4,20]. Linear-time algorithms for computing the diameter, that avoid the computation of the whole distance matrix, have been designed for trees [21], maximal outerplanar graphs [16], interval graphs [25,14], ptolemaic graphs [14], strongly chordal graphs, dually chordal graphs [2], distance-hereditary graphs [11,13] and for graphs of benzenoid systems [4]. It is not yet clear for general graphs whether computing the diameter is easier than computing the whole distance matrix.

In this paper, we study the problem of determining a vertex of high eccentricity of chordal graphs and AT-free graphs. The *eccentricity* of a vertex x is $\text{ecc}(x) = \max_{y \in V} d(x, y)$, where $d(x, y)$ denotes the distance between x and y . The *diameter* of a graph equals the maximum eccentricity achieved by any vertex in the graph. Given v , a vertex of maximum eccentricity, it is trivial to determine the set of vertices whose distance from v equals the diameter of G (these vertices constitute the last layer of a BFS from v).

A graph is *chordal* iff there is no chordless cycle of length more than 3. It is well known that chordal graphs are exactly the intersection graphs of subtrees in trees [3,17]. Interval graphs can be defined as the intersection graphs of subpaths in paths (see [23]). A natural generalization of interval graphs is the concept of directed path graphs. A graph is a *directed path graph* iff it is the intersection graph of a collection of directed paths in a rooted directed tree [18]. A chord $x_i x_j$ in a cycle $C = (x_1, x_2, \dots, x_{2n})$ is an odd chord if, in C , $d(x_i, x_j)$ is odd. A graph is *strongly chordal* [15] if it is chordal and each cycle of even length at least 6 has an odd chord. Strongly chordal graphs represent an interesting subclass of chordal graphs which includes directed path graphs. Three vertices u, v, w are an *asteroidal triple* (AT) if between any two of them, there exists a path that avoids the neighbourhood of the remaining vertex. A graph is *AT-free* if it does not contain an AT. The famous characterization of interval graphs given by Lekkerkerker and Boland says that a graph is interval if and only if it is chordal and AT-free [23].

The algorithm that we present involves two sweeps of the well-known lexicographic breadth first search (LexBFS) introduced by Rose et al. [27] (see algorithm 1) for the recognition of chordal graphs. An example of a LexBFS sweep is presented in Fig. 13. It is somewhat surprising that LexBFS seems to play a fundamental role for both chordal and AT-free graphs, two families that exhibit very little structural similarity (see

for example [8,9,17,26,27]). Dragan et al. [14] and Dragan [12] have proved the following theorem that demonstrates further similar behaviour for chordal and AT-free graphs.

Theorem 1 (Dragan et al. [14], Dragan [12]). *Let v be the vertex visited last by an arbitrary LexBFS. If the graph is chordal or AT-free, then the eccentricity of v is within 1 of the diameter of the graph. If the graph is interval, the eccentricity of v equals the diameter.*

Algorithm 1. Lexicographic breadth first search (LexBFS) [27]

```

Input: A graph  $G = (V, E)$ 
Output: An ordering  $\sigma$  of the vertices of  $V$ 
begin
  assign the label  $\emptyset$  to each vertex ;
  for  $i = n$  to 1 do
    pick an unnumbered vertex  $x$  with the largest label in the lexicographic order ;
    for each unnumbered neighbour  $y$  of  $x$  do
      add  $i$  to  $label(y)$  ;
     $\sigma(i) \leftarrow x$  / * number  $x$  by  $i$  */ ;
  end

```

Note that LexBFS can be started from any vertex of the graph G . We will denote by LexBFS(w) a LexBFS started from vertex w . In this paper, we examine the following very simple 2-sweep LexBFS algorithm and study its performance on chordal and AT-free graphs.

Algorithm 2. 2-sweep LexBFS

```

Input: A graph  $G$ 
Output: A vertex  $v$ 
begin
  Let  $w$  be an arbitrary vertex;
   $u \leftarrow$  the last vertex numbered by LexBFS( $w$ );
   $v \leftarrow$  the last vertex numbered by LexBFS( $u$ );
  return  $v$ ;
end

```

In particular, we examine conditions when $ecc(v) = diam(G) - 1$, where v is the vertex returned by the 2-sweep algorithm. These conditions include forbidden subgraph results for both chordal and AT-free graphs. The forbidden subgraph result for chordal graphs immediately shows that the algorithm works optimally (i.e. $ecc(v) = diam(G)$) for directed path graphs.

Before presenting these results, we show that it is unlikely that the diameter problem on either chordal or AT-free graphs can be solved in quadratic time. To do this we introduce the disjoint sets problem.

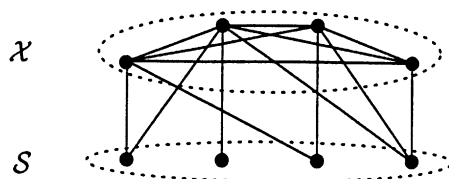


Fig. 1. The set \mathcal{X} is represented by a clique and \mathcal{S} by an independent set. A set S_i is adjacent to its elements in \mathcal{X} . The diameter of this graph is 3 iff there exist two disjoint sets in \mathcal{S} .

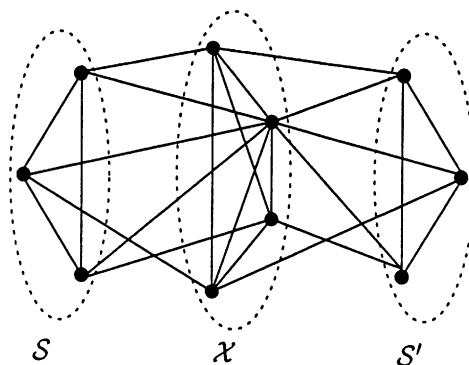


Fig. 2. The set \mathcal{X} is represented by a clique. Two copies of \mathcal{S} are also represented by two cliques. A set S_i (resp. S'_i) is adjacent to its elements in \mathcal{X} . The diameter of this graph is 3 iff there exist two disjoint sets in \mathcal{S} .

2. Disjoint sets problem

Given $\mathcal{S} = \{S_1, S_2, \dots, S_n\}$ sets over the base set \mathcal{X} , the disjoint sets problem (DSP) asks whether there exist i and j such that $S_i \cap S_j = \emptyset$. As pointed out by Chepoi and Dragan in [5], a fast algorithm (i.e. quadratic time or better) for determining whether a split graph (and thus a chordal graph) has diameter 2 or 3 would imply a fast algorithm for the DSP (see Fig. 1).

In Fig. 2, a similar transformation is presented to show that the diameter equals 2 or 3 problem on co-comparability graphs (and thus AT-free graphs) would have the same impact on the DSP.

Thus it seems unlikely that a linear or quadratic time algorithm exists for the diameter problem on either chordal or AT-free graphs. We now present the main results of our paper.

3. Results

The distance between a vertex x and a set of vertices S , denoted by $d(S, x)$, is the minimum distance between x and a vertex of S . The following easy property of LexBFS holds for arbitrary graphs.

Lemma 1. *Let S_i be the numbered vertices at step i of LexBFS. If $x \notin S_i$ and $y \notin S_i$ are two vertices such that $d(S_i, x) < d(S_i, y)$, then x will be numbered before y .*

Let $\sigma = (v_1, v_2, \dots, v_n)$ be an ordering of the vertex set of a graph G . We write $a < b$ whenever in a given ordering σ vertex a has a smaller number than vertex b . Moreover, $\{a_1, \dots, a_l\} < \{b_1, \dots, b_k\}$ is an abbreviation for $a_i < b_j$ ($i = 1, \dots, l$; $j = 1, \dots, k$). An ordering of the vertex set of a graph G generated by LexBFS is called a *LexBFS-ordering*.

In what follows we will often use the following property (cf. [22]):

(P1) If $a < b < c$ and $ac \in E$ and $bc \notin E$ then there exists a vertex d such that $c < d$, $db \in E$ and $da \notin E$.

It is well known that any LexBFS-ordering has property (P1) [19]. Moreover, any ordering fulfilling (P1) can be generated by LexBFS [14].

We now note that the LexBFS algorithm is guaranteed to find the diameter for arbitrary graphs, if the diameter equals 2.

Proposition 1. *Let G be an arbitrary graph and let u be the vertex of G visited last by a LexBFS. If $\text{diam}(G) = 2$, then $\text{ecc}(u) = \text{diam}(G)$.*

Proof. Let (x, y) be a diametrical pair of vertices of G , i.e. $d(x, y) = \text{diam}(G) = 2$. To prove the proposition we just have to show that u is not a universal vertex. Assume both vertices x and y are adjacent to u . Then, since $u < \{x, y\}$, by (P1) there exists a vertex $t > \{x, y\}$ such that $tu \notin E$, thereby proving that u is not universal. \square

3.1. Chordal graphs

We now turn our attention to chordal graphs. An ordering σ of the vertex set of a graph G is a perfect elimination ordering if $bc \in E$ for all vertices, a, b and c with $a < \{b, c\}$ and $ab, ac \in E$. The following theorem presents the well known characterization of chordal graphs.

Theorem 2 (Rose et al. [27]). *Let σ be a LexBFS-ordering of a graph G . Then G is a chordal graph if and only if σ is a perfect elimination ordering of G .*

It is interesting to note that maximum cardinality search (MCS) exhibits the same property as LexBFS with respect to perfect elimination orderings of chordal graphs [29]. For MCS, unlike LexBFS, the eccentricity of the last vertex is not guaranteed to be within any constant of the diameter of the (chordal) graph. (Consider an arbitrarily long path where a new vertex x is adjacent just to a midpoint of the path; it is straightforward to construct an MCS that ends at x .)

Let $P = (x_0 - x_1 - \dots - x_{k-1} - x_k)$ be an arbitrary path of G and let σ be an ordering of the vertex set of this graph. The path P is *monotonic* (with respect to σ)

if $x_0 < x_1 < \dots < x_{k-1} < x_k$ holds whenever $x_0 < x_k$, and P is *convex* if there is an index i ($1 \leq i < k$) such that $x_0 < x_1 < \dots < x_{i-1} < x_i > x_{i+1} > \dots > x_{k-1} > x_k$. Then vertex x_i is called the *switching point* of the convex path P .

In the remainder of this section we assume that G is a chordal graph and σ is a LexBFS-ordering of G .

By Theorem 2 no induced path $P = (x_0 - \dots - x_k)$ of G can contain a vertex x_j ($1 \leq j < k$) with $x_{j-1} > x_j < x_{j+1}$. Hence, we have the following.

Lemma 2. *Every induced path of G is either monotonic or convex.*

Now let $P = (x_0 - \dots - x_k)$ be a shortest path of G connecting x_0 and x_k . We say that P is a *rightmost shortest path* if the sum $x_0 + x_1 + \dots + x_k$ of the positions of x_0, \dots, x_k in σ is largest among all shortest paths connecting x_0 and x_k .

Lemma 3. *Let $P = (x_0 - \dots - x_{2k})$ be a shortest path in G such that the subpath $P' = (x_i - \dots - x_{2k})$, $i \geq k$, of P is a rightmost shortest path connecting x_i and x_{2k} . If $x_0 < x_{2k}$ and x_k is the switching point of P , then $x_{k+j} > x_{k-j}$ holds for each j ($i - k + 1 \leq j \leq k$).*

Proof. We will show that, for each j ($i - k + 1 \leq j \leq k - 1$), if $x_{k+j} < x_{k-j}$ then $x_{k+j+1} < x_{k-j-1}$ holds too. Since $x_{2k} > x_0$ this will give a contradiction.

So, let $x_{k+j} < x_{k-j}$ but $x_{k+j+1} > x_{k-j-1}$. Since P is a convex path and x_k is the switching point of it we have $x_{k+j+1} < x_{k+j}$ and hence $x_{k+j+1} < x_{k-j}$. Applying (P1) to $x_{k-j-1} < x_{k+j+1} < x_{k-j}$ we find a vertex $t > x_{k-j}$ adjacent to x_{k+j+1} and not to x_{k-j-1} . From $x_{k+j+1} < x_{k+j} < \{t, x_{k+j-1}\}$ and Theorem 2 we deduce that t is adjacent to both x_{k+j} and x_{k+j-1} . Then, a contradiction arises to P' being a rightmost shortest path, since $t > x_{k+j}$ and $tx_{k+j+1}, tx_{k+j-1} \in E$. \square

Lemma 4. *Let $P = (x_0 - \dots - x_k)$ be a rightmost shortest path in G which is convex and let x_i be the switching point of P . Then $d(x_0, x_i) \geq d(x_k, x_i)$ whenever $x_0 < x_k$.*

Proof. We prove the assertion by induction on k . Note that any subpath of a rightmost shortest path is again a rightmost shortest path. For $k = 2$ evidently the assertion holds. So, let $k \geq 3$. Since P is convex we have $x_k < x_{k-1}$ and hence $x_0 < x_{k-1}$. By the induction hypothesis, $d(x_0, x_i) \geq d(x_{k-1}, x_i)$. If $d(x_0, x_i) \geq d(x_{k-1}, x_i) + 1$ then $d(x_0, x_i) \geq d(x_k, x_i)$ and we are done. So, assume that $d(x_0, x_i) = d(x_{k-1}, x_i)$. Since $d(x_1, x_i) < d(x_{k-1}, x_i)$ by the induction hypothesis we must have $x_{k-1} < x_1$. Moreover, from $x_k < x_{k-1}$ we conclude $x_k < x_1$. Applying now (P1) to $x_0 < x_k < x_1$ we get a vertex $t > x_1$ adjacent to x_k and not to x_0 . From $x_k < x_{k-1} < \{t, x_{k-2}\}$ and Theorem 2 we deduce that t is adjacent to both x_{k-1} and x_{k-2} . Then, a contradiction arises to P being a rightmost shortest path, since $t > x_{k-1}$ and $tx_k, tx_{k-2} \in E$. \square

Let u be the vertex of a chordal graph G visited last by LexBFS.

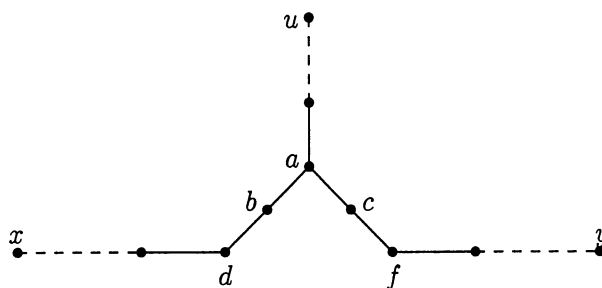


Fig. 3.

Lemma 5. For every two vertices x and y of G such that $d(x, u) \leq d(y, u)$, $d(x, y) \leq d(y, u) + 1$ holds. Moreover, if $d(x, y) = d(y, u) + 1$ then $d(y, u) = d(x, u)$ and $d(y, u)$ is even.

Proof. Assume that $d(x, y) \geq d(y, u) + 1$. Consider in G rightmost shortest paths P_x and P_y , connecting vertex u with vertices x and y , respectively. Let a be the common vertex of the paths P_x and P_y furthest from u . Note that, since a subpath of a rightmost shortest paths is again a rightmost shortest path, paths P_x and P_y coincide in the part from u to a and do not have any other common vertices. Denote the common subpath of these paths by P_a . From $d(x, u) \leq d(y, u)$, we conclude $d(x, a) \leq d(y, a)$.

By Lemma 2, P_x and P_y are monotonic or convex. First we show that these paths cannot have a switching point on the subpath P_a . Assume by way of contradiction that a vertex z of P_a is the switching point of P_x or P_y . Then by Lemma 4 we obtain $d(u, z) \geq d(x, z)$, if z is the switching point of P_x , or $d(u, z) \geq d(y, z)$ if z is the switching point of P_y . Since $d(x, u) \leq d(y, u)$, in both cases we have $d(u, z) \geq d(x, z)$. Hence,

$$\begin{aligned} d(x, y) &\leq d(x, a) + d(a, y) \leq d(x, z) + d(z, y) \\ &\leq d(u, z) + d(z, y) = d(u, y) \leq d(x, y) - 1, \end{aligned}$$

a contradiction.

Now let b and c be the neighbours of a in the paths P_x and P_y , respectively, which do not belong to the monotonic path P_a (see Fig 3). Since $a < \{b, c\}$ by Theorem 2 we get $bc \in E$.

Suppose $c < f$, where f is the neighbour of c in the path P_y distinct from a . If $b > c$ then by Theorem 2 vertices b and f will be adjacent, contradicting P_y being rightmost. Hence, $b < c$ must hold. Now from the fact that the path P_x is rightmost we deduce that $dc \notin E$ and $d < b$, where d is the neighbour of b in the path P_x distinct from a . Since $a < b > d$ the path P_x is convex and b is the switching point of P_x . By Lemma 4 we obtain $d(u, b) \geq d(x, b)$ and hence $d(u, c) \geq d(x, b)$. From $d(x, y) \leq d(x, b) + 1 + d(y, c) \leq d(u, c) + d(y, c) + 1 = d(u, y) + 1$ and our assumption that $d(x, y) \geq d(u, y) + 1$ we see that $d(u, b) = d(x, b)$. Hence from $u < x$ and Lemma 3 it follows that $a < d$. Now we can apply (P1) to $a < d < c$ and get a vertex $t > c$ adjacent to d and not to

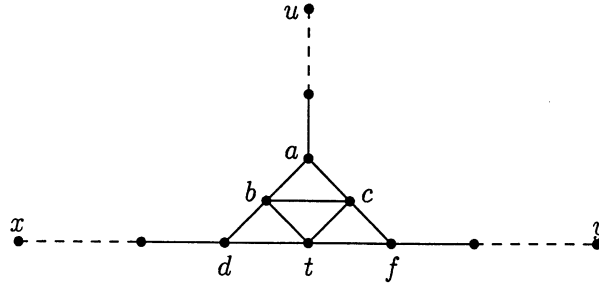


Fig. 4.

a. Since $t > c > b > d$ and $c < f$ by Theorem 2, vertex t must be adjacent to b, c, f (see Fig. 4). Thus,

$$d(x, y) \leq d(x, d) + 2 + d(f, y) = d(x, b) + d(y, c) \leq d(u, c) + d(y, c) = d(u, y),$$

and a contradiction to the assumption $d(x, y) \geq d(u, y) + 1$ arises.

Hence, $c > f$, and therefore P_y is a convex path and c is the switching point of P_y . Again by Lemma 4 we have $d(u, c) \geq d(c, y) \geq d(b, x)$. Hence,

$$d(x, y) \leq d(x, b) + 1 + d(c, y) \leq d(u, c) + 1 + d(c, y) = d(u, y) + 1.$$

Since our assumption was $d(x, y) \geq d(u, y) + 1$ we conclude $d(x, y) = d(u, y) + 1$ and $d(u, c) = d(c, y) = d(b, x)$, i.e. $d(u, y)$ is even and $d(u, x) = d(u, y)$. \square

Theorem 3. Let u be the vertex of a chordal graph G last visited by a LexBFS, and let x, y is a pair of vertices such that $d(x, y) = \text{diam}(G)$. If $\text{ecc}(u) < \text{diam}(G)$ then $\text{ecc}(u)$ is even, $d(u, x) = d(u, y) = \text{ecc}(u)$ and $\text{ecc}(u) = \text{diam}(G) - 1$.

Proof. Assume, without loss of generality, that $d(u, x) \leq d(u, y)$. Then Lemma 5 gives $d(x, y) \leq d(u, y) + 1$. On the other hand, we have $d(x, y) = \text{diam}(G) \geq \text{ecc}(u) + 1 \geq d(u, y) + 1$. Hence, $\text{ecc}(u) = d(u, y)$, $d(x, y) = d(u, y) + 1$ and $\text{diam}(G) = \text{ecc}(u) + 1$. Applying again Lemma 5 we conclude $d(u, x) = d(u, y) = \text{ecc}(u)$ and $\text{ecc}(u)$ is even. \square

We continue with rather surprising results concerning the parity of the diameter of the graph and the parity of the eccentricity of the vertex visited last by LexBFS.

Corollary 1. If the diameter of a chordal graph G is even, then the vertex last visited by a LexBFS has eccentricity equal to $\text{diam}(G)$.

Corollary 2. If the vertex u of a chordal graph G last visited by a LexBFS has odd eccentricity, then $\text{ecc}(u) = \text{diam}(G)$.

To prove the final result on chordal graphs we need the following auxiliary lemmas, the first two of which are well known.

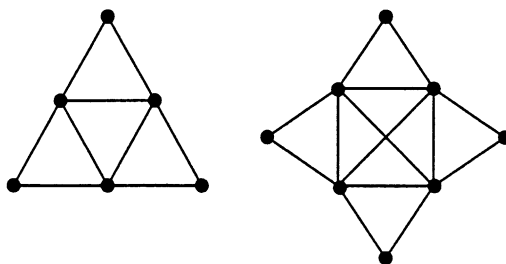


Fig. 5. The 3-sun and the 4-sun.

Lemma 6. For every two vertices x, y of a chordal graph G and every k , $0 < k \leq d(x, y)$, the set $S_k(x, y) = \{z \in V : d(z, x) = k \text{ and } d(z, y) = d(x, y) - k\}$ induces a complete subgraph of G .

Lemma 7. If all vertices of a complete subgraph C of a chordal graph G have the same distance k from a vertex x , then there is a common neighbour z of all vertices of C which is at distance $k - 1$ from x .

Lemma 8. Let G be a chordal graph which does not contain an induced 3-sun (see Fig. 5), let w, x be arbitrary vertices of G and let u be the vertex visited last by $\text{LexBFS}(w)$. Also let P be a rightmost (with respect of $\text{LexBFS}(w)$) shortest path connecting u with x and assume that $d(u, x) = 2k$. Then the vertex z of P with $d(u, z) = k$ lies on a shortest path connecting vertices w and u .

Proof. Consider in G a rightmost shortest path P' connecting vertex u with w and let a be the common vertex of the paths P and P' furthest from u . By Lemma 2, the path P is monotonic or convex (note that the path P' is monotonic since w is the vertex with the largest index in $\text{LexBFS}(w)$). Let b and c be two neighbours of a in the paths P' and P , respectively, which are at distance $d(u, a) + 1$ from u . (The degenerate cases where $a = x$ or $w = a$ are trivial.) We will show that $d(u, a) \geq k - 1$ and if $d(u, a) = k - 1$ then the vertex c belongs to a shortest path connecting w with u .

If $a > c$ then the path P is convex and a is the switching point of P . By Lemma 4 $d(u, a) \geq d(a, x)$. Since $d(u, x) = 2k$ we have $d(u, a) \geq k$.

Assume now that $c > a$. Then from $a < \{b, c\}$ and Theorem 2 we get $bc \in E$. As before, denote by d and f the neighbours of b and c , respectively, on the paths P' and P distinct from a . We have $a < b < d$ and $a < c$. If $b < c$ then by Theorem 2 $dc \in E$ and hence the path P' is not rightmost, a contradiction. So, $b > c$. Analogously, since P is rightmost $c > f$ and $bf \notin E$ must hold. Thus, the path P is convex and c is the switching point of P . From $d(u, x) = 2k$ and Lemma 4 we obtain $d(u, c) \geq k$, i.e. $d(u, a) \geq k - 1$. For the case $d(u, c) = k = d(u, a) + 1$ we will show that the vertex c is adjacent to d . Indeed, if $d(u, c) = k$ then by Lemma 3 $a < f$. Hence, we can apply (P1) to $a < f < b$ and get a vertex $t > b$ adjacent to f and not to a . Since

$t > b > c > f$ and $b < d$ by Theorem 2 vertex t must be adjacent to c, b, d . To avoid an induced 3-sun we must have an edge dc (note that G as a chordal graph does not have any induced cycles). \square

Lemma 9. *Let G be a chordal graph which does not contain an induced 3-sun, w, x, v be arbitrary vertices of G and u be the vertex visited last by LexBFS(w). If $d(u, x) = d(u, v) = d(x, v) = 2k$, then there exists a vertex a of G such that $d(u, a) = d(x, a) = d(v, a) = k$ and $d(w, a) = d(w, u) - k$.*

Proof. Let P_x and P_v be rightmost (with respect to LexBFS(w)) shortest paths connecting u with x and v , respectively. Denote by b and c the midvertices of the paths P_x and P_v , i.e. vertices with $d(x, b) = d(v, c) = k$. By Lemma 8 both vertices b and c belong to shortest paths joining w and u . Hence, if $c = b$ or $d(c, x) = k$ or $d(b, v) = k$ then we are done. So, assume that $c \neq b$ and $d(c, x) > k$, $d(b, v) > k$. Since $d(u, b) = d(u, c) = k$ and $d(w, b) = d(w, c) = d(w, u) - k$, by Lemmas 6 and 7, we get that $bc \in E$ and there are two vertices u', w' in G such that $u'b, u'c, w'b, w'c \in E$ and $d(w', w) = d(w, b) - 1$, $d(u', u) = d(u, b) - 1$. Now consider a cycle of G formed by edge bc and shortest paths connecting b with x , c with v and v with x . Since G is chordal, in this cycle, the middle vertex e of the (v, x) -path must be adjacent to both b and c . Vertices c and e are equidistant from v as well as vertices b and e are equidistant from x . By Lemma 7 there must be vertices x' and v' in G such that $x'b, x'e, v'e, v'c \in E$ and $d(v', v) = d(x', x) = k - 1$. Since $\{d(v, b), d(x, c)\} > k$, $v'b, x'c \notin E$. Hence, to avoid an induced 3-sun, vertex e must be adjacent to both u' and w' . It is easy to see now that e is at distance k from u, x and v and at distance $d(u, w) - k$ from w . \square

The main result of this subsection is the following.

Theorem 4. *If G is a chordal graph and if v , the vertex returned by algorithm 2, is not of maximum eccentricity, then G contains either an induced 3-sun or an induced 4-sun (see Fig. 5) or one of the graphs from Fig. 6 as an induced subgraph.*

Proof. Let w^* be an arbitrary vertex of G , let u^* be the vertex numbered last by LexBFS(w^*) and let v^* be the vertex numbered last by LexBFS(u^*). We will show that if $\text{ecc}(v^*) < \text{diam}(G)$ and G contains neither induced 3-suns nor induced 4-suns then G must contain one of the graphs from Fig. 6 as an induced subgraph.

Let x^*, y^* be a diametral pair of vertices of G , i.e. $d(x^*, y^*) = \text{diam}(G)$. Since $\text{ecc}(v^*) < \text{diam}(G)$, by Theorem 3, we have $\text{ecc}(u^*) = \text{ecc}(v^*) = \text{diam}(G) - 1$, $\text{ecc}(v^*)$ is even, say $\text{ecc}(v^*) = 2k$, and $d(u^*, x^*) = d(u^*, y^*) = d(u^*, v^*) = d(v^*, x^*) = d(v^*, y^*) = 2k = d(x^*, y^*) - 1$. Applying Lemma 9 to x^*, u^*, v^*, w^* and to y^*, u^*, v^*, w^* we will find two vertices a and b such that $d(u^*, a) = d(y^*, a) = d(v^*, a) = d(u^*, b) = d(x^*, b) = d(v^*, b) = k$ and $d(w^*, b) = d(w^*, a) = d(w^*, u^*) - k$. Denote $d(w^*, u^*) - k$ by l . By Lemmas 6 and 7 vertices a and b are adjacent and have three common neighbours u, v and w such that u, v are at distance $k - 1$ from u^*, v^* , respectively, and w is at distance $l - 1$ from

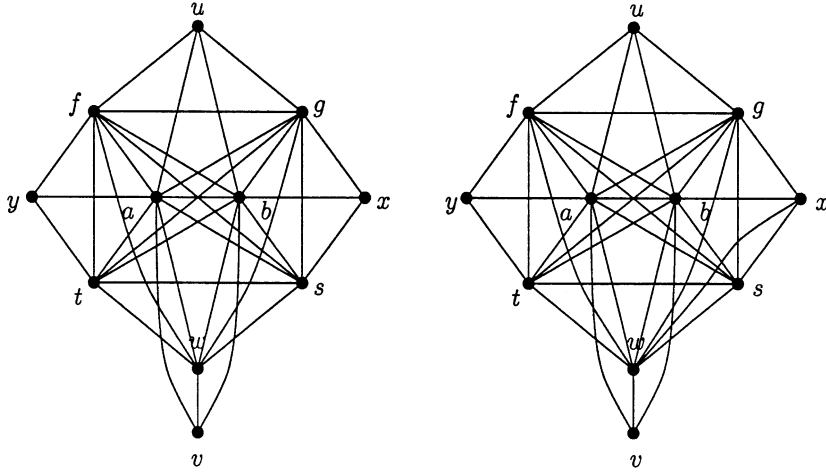


Fig. 6. Strongly chordal graphs where 2 sweeps of LexBFS are not enough to find the diameter. LexBFS(w) and LexBFS(u) orderings of the left graph are $(u, x, y, v, g, f, a, b, s, t, w)$ and $(v, x, y, s, t, w, a, b, g, f, u)$, respectively. Here we have $\text{ecc}(v) = 2 < 3 = d(x, y) = \text{diam}(G)$.

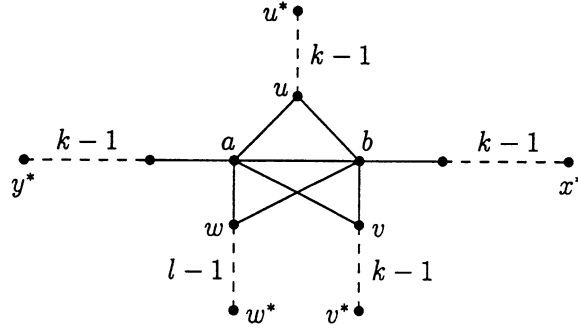


Fig. 7.

w^* (see Fig. 7). We choose v rightmost in LexBFS(w^*), i.e. v has the largest index in LexBFS(w^*) among all vertices which are adjacent to a, b and at distance $k-1$ from v^* .

Let σ be the LexBFS(w^*)-ordering of the vertex set of G with last visited vertex u^* (note that only in the proof of claim 3 do we use the LexBFS(u^*)-ordering). From Lemma 1 and distance requirements we derive $w > \{a, b\} > u$ and $uw, uv \notin E$. Without loss of generality assume that $a < b$.

Claim 1. In the LexBFS(w^*)-ordering, vertex u has a smaller index than vertex v .

Proof. If $a < v$ then we are done since $u < a$. So, assume $a > v$. Consider a rightmost shortest path P_v connecting vertex v with v^* and an arbitrary shortest path P_u joining u with u^* . Let u' and v' be the neighbours of u and v , respectively, on the paths P_u

and P_v . Since $u < a > v$ the shortest path P formed by P_u , P_v and edges ua , av is convex and a is the switching point of P . Consequently, $u' < u$ and $v' < v$. Moreover, from $u^* < v^*$ and Lemma 3 we get $u' < v'$. Now, if $v < u$ then $u' < v' < u$ and we can apply (P1) and find a vertex $z > u$ adjacent to v' and not to u' . Since $v' < v < z$ and $v < a < b$ by Theorem 2 z is adjacent to v , a , b . This gives a contradiction to the choice of v (z is adjacent to both a and b , is at distance $k-1$ from v^* , but $z > u > v$). \square

Claim 2. *There exist two adjacent vertices t and s in G such that $d(x^*, s) = d(y^*, t) = k$ and both t and s are adjacent to a, b, w and not to u . Moreover, $a < s$ and $b < t$ hold in $\text{LexBFS}(w^*)$.*

Proof. Pick the neighbour d of a in a rightmost shortest path connecting a with y^* and the neighbour e of b in a rightmost shortest path connecting b with x^* . From distance requirements we have $d(e, d) = 3$ and $ud, ue \notin E$. Note that w may be adjacent to d or e but only to one of them. Since $db \notin E$ and $a < b$ we get $d < a$. Hence, we can apply Lemma 3 to a shortest path connecting u^* with y^* and passing through vertices u, a, d and get $u < d$. Thus, $u < d < b$, $bu \in E$ and $db \notin E$ hold. By property (P1) there must be a vertex $t > b$ adjacent to d and not to u . Since $d < a < b < t$ and $b < w$, by Theorem 2, t is adjacent to a, b, w . It cannot be adjacent to e because of $d(e, d) = 3$. Consequently, $e < b$ must hold yielding $u < e$ (similar as $u < d$). If now $a < e$ then (P1) applied to $a < e < t$ gives a vertex $p > t$ adjacent to e and not to a . From $e < b < t < p$ and Theorem 2 vertex p is adjacent to b and t . But then vertices u, a, b, d, t and p induce a 3-sun, a contradiction.

Hence, $a > e$. Moreover, every vertex adjacent to both b and t must be adjacent to a as well, otherwise we will have again an induced 3-sun. Now we apply (P1) to $u < e < a$ and get a vertex $s > a$ adjacent to e and not to u . From $e < \{b, s\}$ vertices b and s are adjacent. If $st \notin E$ then $b > s$ holds and (P1) applied to $a < s < t$ gives a vertex $p > t$ adjacent to s and not to a . From $s < b < t < p$ and Theorem 2 vertex p is adjacent to b and t . Since $ap \notin E$ there is a 3-sun on u, a, b, d, t, p . Thus, $st \in E$ and hence $sa \in E$, otherwise there is a 3-sun on s, b, t, a, d, u . From $a < \{w, s\}$, vertex s is adjacent to w too. Finally we note that $d(x^*, s) = d(y^*, t)$ since otherwise $d(x^*, y^*) < 2k + 1$. \square

Claim 3. *There exist two adjacent vertices f and g in G such that $d(x^*, g) = d(y^*, f) = k$ and both f and g are adjacent to a, b, u and not to v .*

Proof. This proof is similar to the proof of claim 2. Instead of the $\text{LexBFS}(w^*)$ -ordering we use the $\text{LexBFS}(u^*)$ -ordering of the vertex set of G with last visited vertex v^* . \square

Now we have all the prerequisites to construct one of the graphs of Fig. 6. Again let σ be a $\text{LexBFS}(w^*)$ -ordering of G with last visited vertex u^* . From the discussion above we have a subgraph of G presented in Fig. 8a.

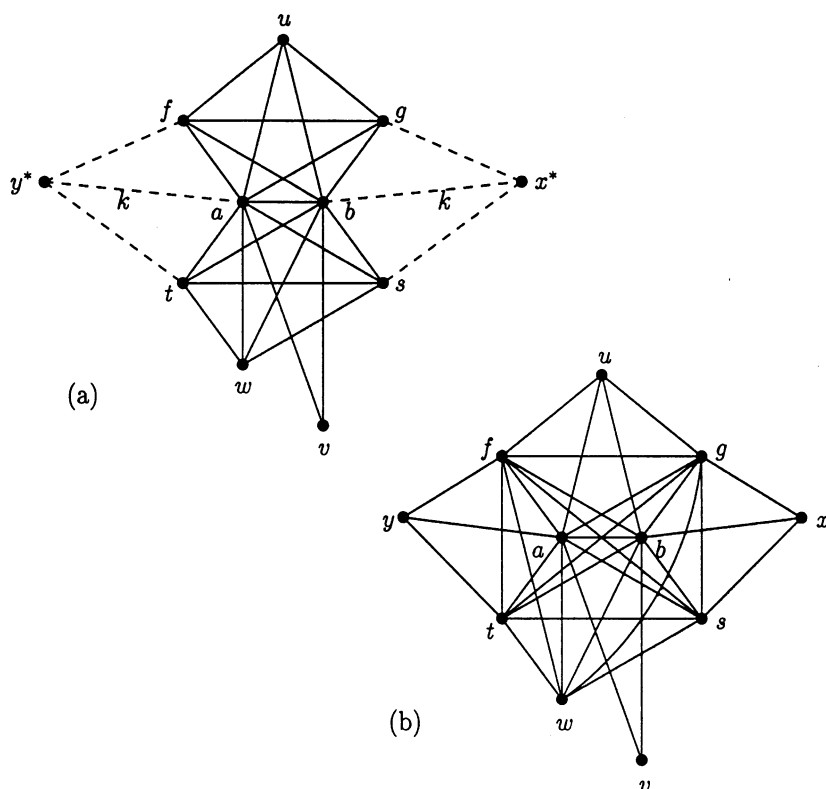


Fig. 8.

Since $b, g, s \in S_k(x^*, y^*)$, by Lemmas 6 and 7 vertices b, g, s are pairwise adjacent and they have a common neighbour x which is at distance $k-1$ from x^* . Analogously, vertices a, f, t are pairwise adjacent and have a common neighbour y at distance $k-1$ from y^* . From distance requirements we have $d(x, y)=3$ and $ux, uy, uw, uv, vx, vy \notin E$. Note that w may be adjacent to x or y but only to one of them. Recall also that $us, ut, vf, vg \notin E$ (see claims 2 and 3). Furthermore, each vertex adjacent to both t and g must be adjacent to f too, otherwise an induced 4-cycle or an induced 3-sun arises. Similarly, each vertex adjacent to both f and s must be adjacent to g as well. Hence, $gt, fs \in E$. If $wf, wg \notin E$ then vertices u, f, g, y, x, t, s, w induce a 4-sun (or a 4-cycle). So, w is adjacent to f and hence to g (or to g and hence to f). Vertex v is adjacent neither to s nor to t . Indeed, if, for example, $vs \in E$ then we get an induced 3-sun formed by v, s, a, g, x, u .

Thus we have constructed a subgraph of G presented in Fig. 8b. Only the following additional edges are possible: wv and/or either wx or wy . If $wv \in E$ then we have a graph from Fig. 6 as an induced subgraph of G . So, assume that $wv \notin E$.

Since vertices a, b, f, g are at distance l (recall that $l = d(w^*, u^*) - k$) from w^* , while vertex w is at distance $l - 1$ and vertex u is at distance $l + 1$, from Lemma 1, $w > \{a, b, f, g\} > u$ must hold. We had also $a < b < t$, $a < s$ and $u < v$ (see claims 1 and 2). From $w > b > a$, $wv \notin E$ and Theorem 2 we derive $v < a < b$. Assume that $f < v$. Then we can apply (P1) to $f < v < w$ and find a vertex $p > w$ adjacent to v and not to f . Again by Theorem 2, p is adjacent to a, b and hence to t, s . Since $pf \notin E$ vertices u, f, g, y, p, t, s, x induce either a 4-sun or a 3-sun (depending on whether p and g are adjacent). Thus, we conclude $v < f$. Now we apply (P1) to $u < v < f$ and get a vertex $p > f$ adjacent to v and not to u . As before p must be adjacent to a, b . We will show that p is adjacent to t, s, f, g as well. If $pt \notin E$ then from $t > b > a$ and Theorem 2 we obtain $p < a$, i.e. $p < a < b < t$ holds. Applying (P1) to $f < p < t$ we find a vertex $q > t$ adjacent to p and not to f . Since $p < a < b < t < q$ and $a < s$, by Theorem 2, q is adjacent to a, b, t, s . But then vertices f, y, u, t, b, q induce a 3-sun or a 4-cycle (if $qy \in E$ or $qu \in E$). Thus, vertex p must be adjacent to t . Analogously we can show that $ps \in E$. Consequently, $pg, pf \in E$ too. Indeed, if $pg \notin E$ then vertices p, s, a, x, g, u induce a 3-sun or a 4-cycle if $(px \in E)$, and if $pf \notin E$ then vertices p, t, b, y, f, u induce a 3-sun. Now it remains to observe that vertices $u, y, x, v, p, a, b, s, t, f$ and g induce a graph from Fig. 6 (replace w with p in those pictures). \square

Recall that a graph G is a *comparability* graph if one can assign directions to edges of G so that the resulting digraph G' is transitive; that is, whenever (x, y) and (y, z) are edges of G' then (x, z) is also an edge of G' [19]. Since none of the graphs from Figs. 5 and 6 is a directed path graph as well as a comparability graph, this theorem immediately yields the following corollary.

Corollary 3. *Algorithm 2 finds a vertex of maximum eccentricity for directed path graphs and for chordal comparability graphs.*

Unfortunately, this result cannot be extended to the strongly chordal graphs since for each of the graphs from Fig. 6, 2 sweeps of LexBFS are not enough to find the diameter. Furthermore, in Fig. 9 we present a strongly chordal graph for which 3 sweeps of LexBFS are also insufficient.

3.2. AT-free graphs

We now turn our attention to AT-free graphs and start by recalling some known results. A pair of vertices (x, y) is said to be a *dominating pair* if for every x, y path P and every vertex $z \in V$, $N(z) \cap P \neq \emptyset$. If $N(z) \cap P = \emptyset$, we say that P *misses* z . For vertices u, v and x , we say that u and v are *unrelated with respect to* x if there is a v, x path that misses u and a u, x path that misses v . As an example of the remarkable similarity exhibited by LexBFS on chordal and AT-free graphs, we note the following characterization of AT-free graphs, a characterization that is analogous to Theorem 2 for chordal graphs.

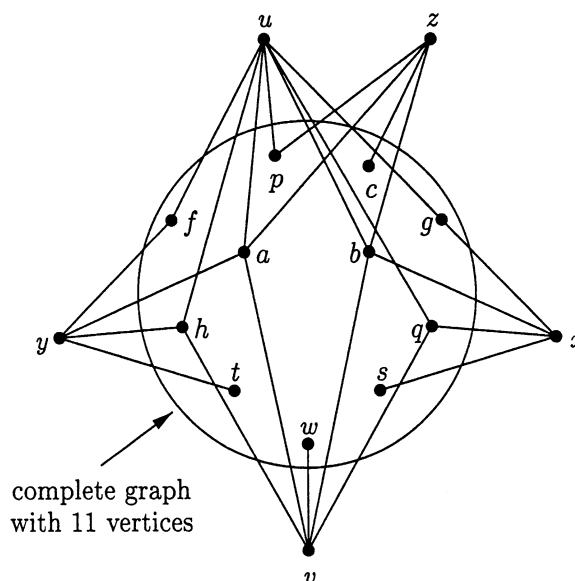


Fig. 9. A strongly chordal graph where 3 sweeps of LexBFS are not enough to find the diameter. LexBFS(w), LexBFS(u), and LexBFS(v) orderings are $(u, z, y, x, v, \dots, c, t, s, w)$, $(v, x, y, z, \dots, g, f, p, u)$ and $(z, u, x, y, \dots, q, h, v)$, respectively. Here we have $\text{ecc}(z) = 2 < 3 = d(x, y) = \text{diam}(G)$.

Theorem 5 (Corneil et al. [8]). *Let σ be a LexBFS ordering of a connected graph G . If G is AT-free, then for all vertices a , b and c with $a < \{b, c\}$, b and c are NOT unrelated with respect to a .*

In [9], it was shown that every connected AT-free graph has a dominating pair; in [8] this was strengthened to show that Algorithm 2 can be used to find such a dominating pair.

Theorem 6 (Corneil et al. [8]). *Let G be a connected AT-free graph and let vertices u and v be as identified in Algorithm 2. Then u and v are a dominating pair of G .*

Furthermore, as the following lemma shows, for sufficiently high diameter, the set of dominating pairs of an AT-free graph can be expressed as the Cartesian product of two distinct sets. (These sets can be found in linear time [8].)

Lemma 10 (Corneil et al. [9]). *If G is a connected AT-free graph with $\text{diam}(G) > 3$ then there exists disjoint vertex sets X, Y such that (x, y) is a dominating pair of G iff $x \in X$ and $y \in Y$.*

The fact that this lemma does not hold for $\text{diam}(G) = 3$ is illustrated by the graph in Fig. 10.

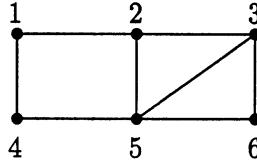


Fig. 10. The dominating pairs are $(1, 3), (1, 6), (4, 3), (4, 6), (1, 5), (2, 5)$ and $(4, 5)$. There are no disjoint sets whose cartesian product defines all dominating pairs.

A weaker version does however hold for AT-free graphs of diameter larger than or equal to 3.

Lemma 11. *Let G be a connected AT-free graph with $\text{diam}(G) \geq 3$ and let V_1 be the set of vertices that are the last vertices of some LexBFS. Then there exists a partition of V_1 into non-empty sets X and Y such that (x, y) is a dominating pair if $x \in X$ and $y \in Y$.*

Proof. If $\text{diam}(G) > 3$, then Lemma 10 applies; X (resp. Y) is the intersection of V_1 with the X (resp. Y) identified in Lemma 10. Thus we only need to consider the case where $\text{diam}(G) = 3$.

Let z be a vertex of maximum eccentricity and let x be a vertex visited last by a LexBFS from z (i.e. $x \in V_1$). We now assume that y is a vertex visited last by a LexBFS from x . Denoting the vertices of distance i from x by $N^i(x)$ we let $X' = V_1 \cap N^3(x)$ and $Y' = V_1 \cap N^3(y)$.

Set $V'_1 = V_1 \setminus (X' \cup Y')$ and let $X'' = V'_1 \cap N^2(x)$ and $Y'' = V'_1 \cap N^2(y)$. Finally we let $X = X' \cup X''$ and $Y = Y' \cup Y''$. We now prove that X and Y satisfy the statement of the lemma.

Claim 4. $X \cap Y = \emptyset$.

Proof. First we show $X' \cap Y' = \emptyset$. Suppose to the contrary that $w \in X' \cap Y'$ and let paths $x - a - b - w$ and $y - c - d - w$ be arbitrary x, w and y, w shortest paths. Trivially, the only possible intersection between these two paths is if $b = d$. Regardless of the intersection, $xc \notin E$ and $ya \notin E$ since otherwise $d(x, y) = 2$. If $b = d$, x and y are unrelated with respect to w contradicting Theorem 5. If $b \neq d$, $xd \notin E$ (else $d(x, w) = 2$) and $yb \notin E$, again showing that x and y are unrelated with respect to w .

$X'' \cap Y'' = \emptyset$ since otherwise the two paths $x - a - w$ and $y - b - w$, where $w \in X'' \cap Y''$, must have $xb \notin E$, $ya \notin E$ and thus x and y would be unrelated with respect to w .

Note that by definition $X' \cap Y' = X'' \cap Y'' = \emptyset$. \square

Claim 5. $X \cup Y = V_1$.

Proof. If $w \in V_1 \setminus (X \cup Y)$, then since $\text{diam}(G) = 3$, $w \in \{x\} \cup N^1(x)$ and $w \in \{y\} \cup N^1(y)$ which implies that $d(x, y) \leq 2$ contradicting $d(x, y) = 3$. \square

Before proceeding to show that every element of X forms a dominating pair with every element of Y , we note the following:

Let u be an arbitrary element of X and v an arbitrary element of Y . Since $u \in X$, $d(x, u) = 2$ or 3 . If $d(x, u) = 3$, $d(y, u) \leq 2$ (by claim 4). If $d(x, u) = 2$, $d(y, u) \leq 1$ (if $d(y, u) = 3$, $u \in Y$ contradicting claim 4; $d(y, u) \neq 2$ since $X'' \cap Y'' = \emptyset$). Similarly if $d(y, v) = 3$, $d(x, v) \leq 2$ and if $d(y, v) = 2$, $d(x, v) \leq 1$.

Claim 6. For all $u \in X$, $v \in Y$, u, v is a dominating pair.

Proof. Suppose to the contrary, that there is a pair of vertices u, v and an induced path P joining them that misses vertex t . We now do a case by case analysis based on the distance from x to P and from y to P . We let $d(x, P)$ denote the length of a shortest path from x to P . Thus $d(x, P) = 0$ means x is on P ; $d(x, P) = 1$ means x has a neighbour on P and $d(x, P) > 1$ means x misses P (i.e. x plays the role of t).

Case 1: $d(x, P) = 0$ and $d(y, P) = 0$ (i.e. both x, y are on P).

This is clearly impossible since the subpath of P between x and y misses t , contradicting x, y being a dominating pair.

Case 2: $d(x, P) = 0$ and $d(y, P) = 1$ (i.e. one vertex (without loss of generality x) is on P , the other is adjacent to P)

First we note (and this also applies to case 3 below) that x must be one of the endpoints of P or adjacent to one of the endpoints. Otherwise u and v are unrelated with respect to x . Since $d(x, u) = 2$ or 3 , either $xv \in E$ or $x = v$. Let p be the neighbour of x on P in the direction towards u . Since $d(x, u) \geq 2$, $u \neq p$.

Since $d(y, P) = 1$, there is a vertex $q \in P$ such that $yq \in E$. Clearly such a vertex q is between u and p but cannot be p (otherwise $d(x, y) = 2$). The x, y path $x \sim (P) \sim q - y$ (this notation indicates the subpath of P from x to q together with the edge qy) shows that $ty \in E$ (otherwise x, y is not a dominating pair). Now let q be the neighbour of y on P that is closest to u . The paths $u \sim (P) \sim q - y - t$ and $u \sim (P) \sim v$ show that v and t are unrelated with respect to u .

Case 3: $d(x, P) = 0$ and $d(y, P) > 1$ (i.e. one vertex (without loss of generality x) is on P , the other is not adjacent to P).

Note that y can now play the role of t . As in case 2 we let p denote the neighbour of x on P and see that $py \notin E$ (otherwise $d(x, y) = 2$). Since $d(y, u) \leq 2$ and $d(y, P) > 1$ we see that $d(y, u) = 2$ and let $u - \alpha - y$ denote such a path ($\alpha \notin P$).

We immediately see that $x\alpha \notin E$ (otherwise $d(x, y) = 2$) and thus paths $u - \alpha - y$ and $u \sim (P) \sim x$ show that x, y are unrelated with respect to u .

Case 4: $d(x, P) = 1$ and $d(y, P) = 1$ (i.e. neither x nor y is on P but each intercepts P).

The x, y path that has all internal vertices on P shows that $xt \in E$ or $yt \in E$ but not both. Without loss of generality assume $xt \in E$. If $uv \in E$, then $yu, xv \in E$. Now y, t are unrelated with respect to v . Thus $uv \notin E$. Let x' be the neighbour of x (on P) that is closest to v and consider the paths:

$v \sim (P) \sim u$ and $v \sim (P) \sim x' - x - t$ (where possibly $v = x'$). Since $xu \notin E$, t and u are unrelated with respect to v unless x' is adjacent to u . Thus $d(x, u) = 2$ implying that $d(y, u) \leq 1$. Since y is not on P , $d(y, u) = 1$. Also $x' \neq v$ since $uv \notin E$.

Now the paths $v \sim (P) \sim u - y$ and $v \sim (P) \sim x' - x - t$ show that y and t are unrelated with respect to v unless y is adjacent to some vertex on the v, x' subpath of P . Let y' be the neighbour of y closest to v ; $y' \neq v$ since $d(y, v) \geq 2$. The existence of y' shows that $x'v \notin E$. Now the path $y - u - x' - x$ misses v contradicting x, y being a dominating pair. Note $xv \notin E$ since x' is the neighbour of x that is closest to v .

Case 5: $d(x, P) = 1$ and $d(y, P) > 1$ (i.e. x intercepts P but y misses P ; note y now plays the role of t).

Since $d(y, P) > 1$ and $d(y, u) \leq 2$, $d(y, u) = 2$ by the path $u - \alpha - y$ where $\alpha \notin P$. Now $xy, yx', x\alpha \notin E$ and also $ux \notin E$. But paths $u - \alpha - y$ and $u \sim (P) \sim x' - x$ show that x and y are unrelated with respect to u .

Case 6: $d(x, P) > 1$ and $d(y, P) > 1$ (i.e. both x and y miss P).

Now $d(y, u) = 2$ by the path $u - \alpha - y$ and $d(x, v) = 2$ by the path $v - \beta - x$ where $\alpha, \beta \notin P$ ($\alpha \neq \beta$). Again $yx, y\beta, x\alpha \notin E$. But path $u - \alpha - y$ and the path induced on $u \sim (P) \sim v - \beta - x$ show that x and y are unrelated with respect to u .

This completes the proof of the lemma. \square

For the graph shown in Fig. 10, $X = \{6\}$ and $Y = \{1, 4\}$. The next proposition presents further facts about the structure of the AT-free graphs.

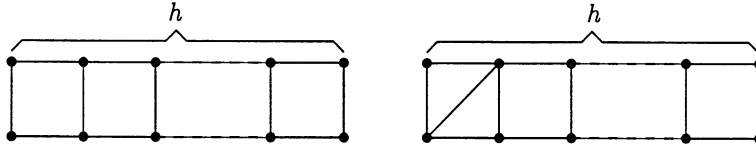
Proposition 2. *Let G be an AT-free graph with $\text{diam}(G) = k > 2$. If $\text{ecc}(v) = k - 1$, where v is the vertex returned by algorithm 2, and u', v' achieve the diameter where $d(u, u') \leq d(u, v')$ then:*

1. $d(u, v) = d(u, v') = d(u', v) = k - 1$,
2. $uu' \in E$ and $vv' \in E$.

Proof. 1. By Lemma 11, u, v is a dominating pair. Thus each of u', v' is adjacent to at least one vertex of Q , an arbitrary shortest u, v path. Let a be the furthest vertex from u (on Q) that is adjacent to u' and let b be the furthest vertex from v (on Q) that is adjacent to v' . By Theorem 1, it is clear that $\text{ecc}(u) = k - 1$ and $d(u, v) = k - 1$. Since $d(u', v') = k$, either $a = u$ or $b = v$ (or both). Without loss of generality assume $v = b$. Now $d(u', v) = k - 1$ since otherwise $d(u', v') < k$. If $u'u \in E$ then the same argument shows that $d(u, v') = k - 1$. Thus we may assume that $uu' \notin E$ and a is the neighbour of u on Q . Suppose $d(u, v') < k - 1$ as witnessed by path R . Now u' must be adjacent to some vertex of R since otherwise the u, v path consisting of R plus the edge $v'v$ misses u' contradicting u, v being a dominating pair. But now, $d(u', v') < k$.

2. This follows immediately from the preceding argument. \square

Although even the 2-sweep LexBFS algorithm does not guarantee a maximum eccentricity vertex for AT-free graphs (as well as for chordal graphs), the previous

Fig. 11. The h -ladder and the h -*-ladder.

proposition shows that such a vertex is in the last BFS layer from v , the vertex returned by algorithm 2.

Before presenting the final result on AT-free graphs, we introduce the notion of an h -ladder and an h -*-ladder.

Definition 1. An h -ladder consists of a chain of h 4-cycles where the 4-cycles are attached as shown in Fig. 11. In an h -*-ladder the first 4-cycle has a diagonal.

Theorem 7. If G is an AT-free graph with $\text{diam}(G) = k \geq 3$ and $\text{ecc}(v) = k - 1$, where v is the vertex returned by algorithm 2, then G contains an induced $(k - 1)$ -ladder or an induced $(k - 1)$ -*-ladder.

Proof. Let u', v' be a diametrical dominating pair such that $u, u' \in X$, $v, v' \in Y$. Clearly such u', v' exist (i.e. $u', v' \in V_1$). Furthermore let P be an arbitrary rightmost shortest u', v path $u' = u_0 - u_1 - \dots - u_{k-1} = v$. Throughout the proof $a < b$ indicates that vertex a has a smaller number than vertex b in the last LexBFS of Algorithm 2.

$L_i (0 \leq i \leq k - 1)$ will denote the BFS layers of G with respect to u (i.e. $L_0 = \{u\}$; $v \in L_{k-1}$ and it has been shown in Proposition 2 that $u' \in L_1, v' \in L_{k-1}$). The following facts follow immediately.

Fact 1. For all $0 \leq i < k - 1$, no neighbour of u_i may be of distance $\leq k - i - 2$ from v' (otherwise $d(u', v') < k$.)

Fact 2. For exactly one $l, 1 \leq l \leq k - 1$, L_l contains 2 vertices from P namely u_{l-1}, u_l . All other layers contain exactly one vertex from P . (Pigeonhole principle and distance properties.) If L_l contains u_{l-1}, u_l we say that P “jogs” in L_l .

Claim 7. If P has its jog in $L_l, 1 \leq l \leq k - 1$, then there is an induced $(k - l - 1)$ -ladder on $\{u_l, \dots, u_{k-1} = v\} \cup \{\alpha_l, \dots, \alpha_{k-1} = v'\}$ where $\alpha_j \in L_j$ for all $j, l \leq j \leq k - 1$ and $\alpha_j > u_j$.

Proof. By induction on $j: k - 1, \dots, l$. If $j = k - 1, \alpha_{k-1} = v', v'v \in E$ (i.e. we have a $[k - (k - 1) - 1 = 0]$ -ladder), $v' > v$ and $v' \in L_{k-1}$.

Now assume the claim is true for $j > l$ and show it's true for $j - 1$. $u_{j-1} \in L_{j-1}$ and thus $u_{j-1} > \alpha_j$. By fact 1, $u_{j-1}\alpha_j \notin E$. Since $\alpha_j > u_j$, by (P1) there exists $\alpha_{j-1} > u_{j-1}$ such that $\alpha_{j-1}\alpha_j \in E, \alpha_{j-1}u_j \notin E$. By Lemma 11, P is a dominating path and thus must

dominate α_{j-1} . By fact 1, $u_{j-2}\alpha_{j-1} \notin E$ and thus $\alpha_{j-1}u_{j-1} \in E$, thereby extending the ladder. \square

As an immediate corollary of claim 7, we see that if $l = 1$ then u is adjacent to $u_0(=u'), u_1$ and α_1 , thereby resulting in a $(k-1)$ -ladder.

Henceforth we assume that $l > 1$. We now show that there is only one possible ordering of u_l, u_{l-1}, α_l .

Claim 8. $u_{l-1} > \alpha_l > u_l$.

Proof. By claim 7, $\alpha_l > u_l$. If the stated order is not present then $\alpha_l > u_{l-1}$. Since $l > 1$, u_{l-2} exists, in L_{l-1} . By fact 1, $u_{l-2}\alpha_l \notin E$ and thus by (P1) there exists $\alpha_{l-1} > u_{l-2}$ such that $\alpha_{l-1}u_{l-1} \notin E$, $\alpha_{l-1}\alpha_l \in E$. By fact 1, $\alpha_{l-1}u_{l-2} \notin E$. Now consider any direct u, α_{l-1} path and append to it the path $\alpha_{l-1} - \alpha_l - \dots - v'$. This path (from u to v') misses u_{l-1} contradicting Lemma 11. \square

Since $\alpha_l > u_l$, there is α_{l-1} (with as large a number in σ as possible) such that $\alpha_{l-1} > u_{l-1}, \alpha_{l-1}\alpha_l \in E$, $\alpha_{l-1}u_l \notin E$. Furthermore, by fact 1, $\alpha_{l-1}u_{l-2} \notin E$. Suppose $\alpha_{l-1}u_{l-1} \notin E$. If $\alpha_{l-1} \in L_{l-1}$, then P misses α_{l-1} (note in this case, $\alpha_{l-1}u_{l+1} \notin E$ since $u_{l+1} \in L_{l+1}$). Thus $\alpha_{l-1} \in L_l$ and $u_{l-2} > \alpha_{l-1}$. Thus there exists $\alpha_{l-2} > u_{l-2}$ such that $\alpha_{l-2}\alpha_{l-1} \in E$, $\alpha_{l-2}u_{l-1} \notin E$. But now a direct path from u to α_{l-2} concatenated with the path $\alpha_{l-2} - \alpha_{l-1} - \dots - v'$ yields a path from u to v' that misses u_{l-1} , contradicting Lemma 11. Thus, $\alpha_{l-1}u_{l-1} \in E$. Furthermore $\alpha_{l-1}u_{l+1} \notin E$ since otherwise the path P' formed from P by replacing u_l with α_{l-1} would contradict P being rightmost. We now examine the possible relative orders of u_{l-2} and α_{l-1} .

Case 1: $\alpha_{l-1} > u_{l-2}$ (i.e. $\alpha_{l-1} \in L_{l-1}$).

If $l = 2$, u is adjacent to α_{l-1} and u_{l-2} and we have an induced $(k-1)$ -ladder. If $l > 2$, u_{l-3} exists (in L_{l-2}) with $u_{l-3}u_{l-2} \in E$, $u_{l-3}\alpha_{l-1} \notin E$ (fact 1). Since $\alpha_{l-1} > u_{l-2}$, there exists $\alpha_{l-2} > u_{l-3}$ such that $\alpha_{l-2}\alpha_{l-1} \in E$, $\alpha_{l-2}u_{l-2} \notin E$. But now a direct u, α_{l-2} path concatenated with the path $\alpha_{l-2} - \alpha_{l-1} - \dots - v'$ yields a path from u to v' that misses u_{l-2} , contradicting Lemma 11.

Case 2: $u_{l-2} > \alpha_{l-1}$.

We differentiate on whether $\alpha_{l-1} \in L_l$ or L_{l-1} .

Case 2.1: $\alpha_{l-1} \in L_l$.

Since $\alpha_{l-1} > u_{l-1}$, there is $\alpha_{l-2} > u_{l-2} (\alpha_{l-2} \in L_{l-1})$ such that $\alpha_{l-2}\alpha_{l-1} \in E$, and $\alpha_{l-2}u_{l-1} \notin E$. First we see that $\alpha_{l-2}\alpha_l \notin E$ since otherwise a direct path from u to α_{l-2} concatenated with the path $\alpha_{l-2} - \alpha_l - \dots - v'$ yields a path from u to v' that misses u_{l-1} , contradicting Lemma 11. Secondly $\alpha_{l-2}u_l \notin E$ since otherwise there would exist vertex $\beta > \alpha_{l-2}$ such that $\beta\alpha_l \in E$, $\beta u_l \notin E$ but this would contradict α_{l-1} having as large a number in σ as possible. Now we see that $\alpha_{l-2}u_{l-2} \in E$. Otherwise, if $l = 2$, P misses α_{l-2} . If $l > 2$, there exists $\alpha_{l-3} > u_{l-3}$ such that $\alpha_{l-3}\alpha_{l-2} \in E$, $\alpha_{l-3}u_{l-2} \notin E$ but now a direct path from u to α_{l-3} concatenated with the path $\alpha_{l-3} - \alpha_{l-2} - \dots - v'$ yields a path from u to v' that misses u_{l-2} , contradicting Lemma 11.

If $l = 2$, $u' = u_{l-2}$ and there is a $(k-1)$ -ladder on $P \cup \{\alpha\} \cup v'$. If $l > 2$, a straightforward induction argument shows that for all i , $1 \leq i \leq l-1$, L_{l-i} contains vertices $\alpha_{l-i-1} > u_{l-i-1}$ where $\alpha_{l-i-1}\alpha_{l-i} \in E$, $\alpha_{l-i-1}u_{l-i-1} \in E$, $\alpha_{l-i}u_{l-i-1} \notin E$ and $\alpha_{l-i-1}u_{l-i} \notin E$. When $i = l-1$, $u_{l-i-1} = u'$ and we have an induced $(k-1)$ -ladder on $P \cup \{\alpha\} \cup v'$.

Case 2.2: $\alpha_{l-1} \in L_{l-1}$.

If $l = 2$ (i.e. $u_{l-2} = u'$), then there is a $(k-1)$ -ladder on $P \cup u \cup \{\alpha\} \cup v'$.

If $l > 2$, let α_{l-2} be the neighbour of α_{l-1} in L_{l-2} with the largest number. Since $u_{l-1} \in L_l$, $\alpha_{l-2}u_{l-1} \notin E$. $\alpha_{l-2}u_{l-2} \in E$ since otherwise P would miss α_{l-2} . Now consider $u_{l-3} \in L_{l-2}$. If $l=3$, $u_{l-3} = u'$ and we have a $(k-1)$ -ladder on $P \cup u \cup \{\alpha\} \cup v'$. If $l > 3$ we now show that $u_{l-3} > \alpha_{l-2}$. Assume to the contrary. Since $l > 3$, u_{l-4} exists (in L_{l-3}). By fact 1, $u_{l-4}\alpha_{l-2} \notin E$. $\alpha_{l-2} > u_{l-3}$ implies there exists $\alpha_{l-3} > u_{l-4}$ such that $\alpha_{l-3}\alpha_{l-2} \in E$, $\alpha_{l-3}u_{l-3} \notin E$ and $u_{l-4}\alpha_{l-3} \notin E$ (by fact 1). Thus P misses α_{l-3} .

Now a straightforward induction argument shows that for all i , $2 \leq i \leq l-1$, L_{l-i} contains vertices $u_{l-i-1} > \alpha_{l-i}$ where $\alpha_{l-i}\alpha_{l-i+1} \in E$, $\alpha_{l-i}u_{l-i} \in E$ (otherwise P misses α_{l-i}), $\alpha_{l-i}u_{l-i-1} \notin E$ (fact 1). When $i = l-1$, $u_{l-i-1} = u'$ and we have a $(k-1)$ -ladder on $P \cup u \cup \{\alpha\} \cup v'$.

This completes the proof of the theorem. \square

Note that this theorem considerably strengthens the following result by Dragan [12]. An *HHD-free graph* does not contain an induced house (complement of P_5) or an induced hole (an odd cycle of length at least 5) or an induced domino (a 2-ladder).

Theorem 8 (Dragan [12]). *If G is an HHD-free, AT-free graph, then the vertex visited last by a LexBFS has maximum eccentricity.*

Similarly considering $K_{1,3}$ induced subgraphs we have:

Corollary 4. *If G is an AT-free graph with no $K_{1,3}$ then the vertex returned by algorithm 2 has maximum eccentricity.*

4. Concluding remarks

First of all, the reader should note a kind of duality in the results when algorithm 2 finds a vertex whose eccentricity is not maximum. For chordal graphs, each of the forbidden subgraphs has an AT. For AT-free graphs, the h -ladder and the h^* -ladder are built with 4-cycles, the smallest non-chordal graph.

Having seen the power of the 2-sweep LexBFS algorithm, it is natural to ask whether significant improvements can be achieved by performing c sweeps for some $c > 2$. In particular, can we find a vertex of maximum eccentricity, although in light of the results of Section 2, this is highly unlikely for c a constant? As shown by the graphs

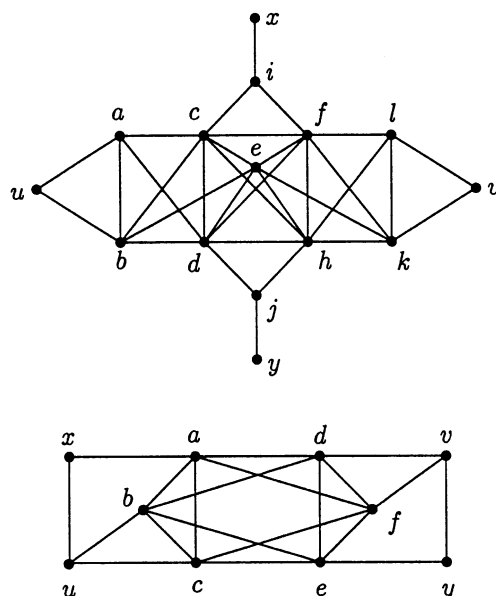
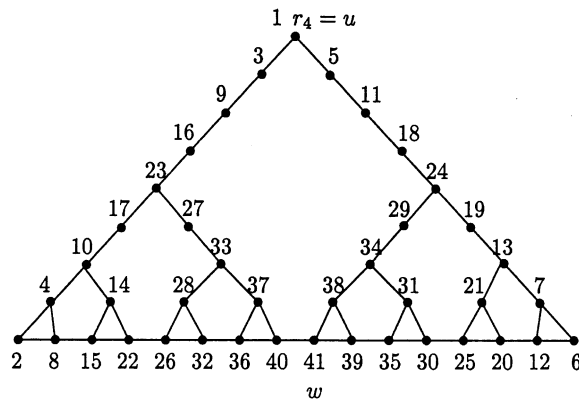
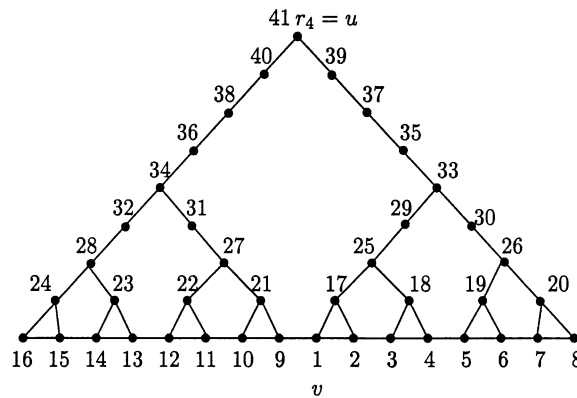


Fig. 12. A chordal graph and an AT-free graph where an infinite number of LexBFS sweeps never end at a maximum eccentricity vertex.

in Fig. 12, for no c , is the c -sweep algorithm guaranteed to find a vertex of maximum eccentricity. The first graph is chordal, the second AT-free. In both graphs any LexBFS starting at u must end at v and vice versa. Thus if the initial choice of vertex is either u or v , a multi-sweep LexBFS algorithm will forever alternate between u and v , thereby missing x and y , the two vertices of maximum eccentricity.

A second obvious question concerns the power of the 2-sweep algorithm on arbitrary graphs. Unfortunately, the answer again is negative. In particular, for any $i > 1$, there is a graph G_i where $\text{ecc}(v) = \text{diam}(G_i) - 2^{i-1} + 1$, where v is the vertex returned by algorithm 2. We construct G_i as follows: Let T_1 be a 2-leaf tree with root r_1 . T_i , $i > 1$, is formed from two copies of T_{i-1} by making r_i , the root of T_i , adjacent to the two r_{i-1} roots. Each $r_i r_{i-1}$ edge then has $2^{i-2} - 1$ new vertices inserted. Finally G_i is formed from T_i by creating a path on the leaves of T_i in the obvious way. G_4 is shown in Fig. 13. If w is the leftmost leaf of the right T_{i-1} and the next vertex chosen in the LexBFS from w is the rightmost leaf of the left T_{i-1} , then the LexBFS will end at $u = r_i$ (see Fig. 13). If the second LexBFS starts at u and breaks ties by choosing the last eligible vertex in the previous sweep, then v , the last vertex, is the same as w (see Fig. 14). It is easy to see that $\text{ecc}(v) = 2^{i-1}$ and $\text{diam}(G_i) = 2^i - 1$ as witnessed by the extreme leaves.

As a final comment, we note that the results in this paper add to the growing evidence of the similar roles played by LexBFS for chordal and AT-free graphs. It would be interesting to find a structural result to explain this surprising phenomenon.

Fig. 13. G_4 together with the first LexBFS.Fig. 14. G_4 and the second LexBFS.

Acknowledgements

D.C. wishes to thank NSERC for financial assistance and LIRMM for their hospitality during his visit, and F.D.'s research was supported by the DFG. The authors would like to thank the anonymous referees for their helpful comments.

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